

Algebraic Complexity Theory

SS 2014, Exercise Sheet #2

$$\begin{pmatrix} p_0 \cdot q_0 \\ p_1 \cdot q_0 + p_0 \cdot q_1 \\ p_2 \cdot q_0 + p_1 \cdot q_1 + p_0 \cdot q_2 \\ p_2 \cdot q_1 + p_1 \cdot q_2 \\ p_2 \cdot q_2 \end{pmatrix} \stackrel{(*)}{=} \begin{pmatrix} 1 & x_0 & x_0^2 & x_0^3 & x_0^4 \\ 1 & x_1 & x_1^2 & x_1^3 & x_1^4 \\ 1 & x_2 & x_2^2 & x_2^3 & x_2^4 \\ 1 & x_3 & x_3^2 & x_3^3 & x_3^4 \\ 1 & x_4 & x_4^2 & x_4^3 & x_4^4 \end{pmatrix}^{-1} \cdot \begin{pmatrix} (p_0 + p_1 \cdot x_0 + p_2 \cdot x_0^2) \cdot (q_0 + q_1 \cdot x_0 + q_2 \cdot x_0^2) \\ (p_0 + p_1 \cdot x_1 + p_2 \cdot x_1^2) \cdot (q_0 + q_1 \cdot x_1 + q_2 \cdot x_1^2) \\ (p_0 + p_1 \cdot x_2 + p_2 \cdot x_2^2) \cdot (q_0 + q_1 \cdot x_2 + q_2 \cdot x_2^2) \\ (p_0 + p_1 \cdot x_3 + p_2 \cdot x_3^2) \cdot (q_0 + q_1 \cdot x_3 + q_2 \cdot x_3^2) \\ (p_0 + p_1 \cdot x_4 + p_2 \cdot x_4^2) \cdot (q_0 + q_1 \cdot x_4 + q_2 \cdot x_4^2) \end{pmatrix}$$

EXERCISE 3:

Recall Karatsuba's Algorithm for polynomial multiplication using $\mathcal{O}(n^{\log_2 3}) \subseteq \mathcal{O}(n^{1.585})$ arithmetic operations. Now fix an algebra \mathcal{A} over the infinite field \mathbb{F} .

- Verify the above identity (*) for any pairwise distinct $x_0, x_1, \dots, x_d \in \mathbb{F}$ and arbitrary $p_0 + p_1 \cdot X + p_2 \cdot X^2, q_0 + q_1 \cdot X + q_2 \cdot X^2 \in \mathcal{A}[X]$.
- Choose $x_j = j$, say, and conclude that two quadratic polynomials over \mathcal{A} can be multiplied using 5 — instead of 9 — multiplications in \mathcal{A} (and arbitrary many additions in \mathcal{A} as well as multiplications by constants from \mathbb{F}).
- Derive an algorithm for multiplying two polynomials over \mathcal{A} of degree n using $\mathcal{O}(n^{\log_3 5}) \subseteq \mathcal{O}(n^{1.465})$ arithmetic operations and constants from \mathbb{F} .
- Generalize a) and b) in order to obtain an algorithm multiplying $p \in \mathcal{A}[X]$ of $\deg(p) \leq k$ and $q \in \mathcal{A}[X]$ of $\deg(q) \leq \ell$ using $k + \ell + 1$ multiplications in \mathcal{A} (and arbitrary many additions in \mathcal{A} as well as multiplications by constants from \mathbb{F}).
Can you identify Karatsuba as a special case?
- Derive, for any fixed $\varepsilon > 0$, an algorithm multiplying two polynomials over \mathcal{A} of degree at most n using $\mathcal{O}(n^{1+\varepsilon})$ arithmetic operations and constants from \mathbb{F} .

EXERCISE 4:

Formalize the following algorithms as straight-line programs and analyze their costs:

- Compute the determinant of a given 3×3 -matrix using Sarrus' Rule.
- Compute the determinant of a given $n \times n$ -matrix using Laplace's Expansion.
- Compute the determinant of a given $n \times n$ -matrix using Leibnitz' Formula.
- Compute the determinant of a given 3×3 -matrix via its LU-decomposition/Gaussian Elimination.