

Advanced Complexity Theorie

SS 2011, Exercise Sheet #1

EXERCISE 1:

Let $L := \{ \vec{x} 0^{|\vec{x}|} \vec{x} : \vec{x} \in \Sigma^* \}$.

- a) Describe a 1-tape Turing machine deciding L in time $\mathcal{O}(n^2)$.
- b) Describe a 2-tape Turing machine deciding L in time $\mathcal{O}(n)$.

EXERCISE 2:

- a) Describe a sequence \vec{x}_n of binary strings of length $|\vec{x}_n| = n$ with $K(\vec{x}_n) \leq \mathcal{O}(\log |\vec{x}_n|)$.
How about \vec{x} with $K(\vec{x}) \leq \mathcal{O}(\log \log |\vec{x}|)$?
- b) Prove that, asymptotically, ‘almost every’ binary string of length n has Kolmogorov Complexity at least $n/2$. Determine the complexity of a random binary string of length n as $n \rightarrow \infty$.
- c) It is not known whether transcendental number π is **normal** (with respect to base 2, say).
Let $\vec{x}_n = (1, 1, 0, 0, 1, 0, 0, 1, 0, 0, 0, 0, 1, 1, 1, 1, 1, 0, 1, 1, 0, 1, 0, \dots) \in \{0, 1\}^n$ denote the sequence of the first n binary digits of π . Estimate the Kolmogorov Complexity of \vec{x}_n
 - i) in case π is normal, and
 - ii) in case it is not.
- d) Let $r \in [0, 1)$ denote an arbitrary real number and \vec{x}_n the sequence of the first n binary digits of r and \vec{y}_n the sequence of the first n *decimal* digits of r , each separately encoded in binary.
Compare the Kolmogorov Complexity of \vec{x}_n with that of \vec{y}_n up to an additive constant independent of n .