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Introduction to Mathematical Software

7th Exercise Sheet

Exercise 1 (Fibonacci Numbers)

- (a) Write a routine int fib(int n) that computes the n-th Fibonacci number by recursively calling itself.
- (b) Now write an iterative routine for computing the n th Fibonacci using a forloop.
- (c) Test both algorithms with n = 42. What do you observe? Does this mean recursion as such is slow?
- (d) What happens if you choose n = 50?

Exercise 2 (Matrix Multiplication)

Write a program to multiply two matrices.

More specifically, you are given an $l \times m$ -matrix A and an $m \times n$ -matrix B. Compute the $l \times n$ -matrix C with $C = A \cdot B$.

Hints:

- use the **#define** preprocessor directive to set the dimensions of the matrices, e.g. **#define L 3**
 - #define M 2
 - #define N 4

L, M, N can then be used like ordinary variables of type int.

• write a subfunction void MatMatMult (double A[L][M], double B[M][N], double C[L][N]) that computes AB and stores the result in C.

Exercise 3 (Efficient Fibonacci)

An efficient way for computing large Fibonacci numbers is the matrix multiplication algorithm. This algorithm is based on the following observations:

The Fibonacci sequence can be described as a 2-dimensional system of difference equations

$$\begin{pmatrix} F_{n+2} \\ F_{n+1} \end{pmatrix} = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix} \begin{pmatrix} F_{n+1} \\ F_n \end{pmatrix}$$

where F_n denotes the n-th Fibonacci number.

This yields the following closed form for the computation of Fibonacci numbers

$$\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^n = \begin{pmatrix} F_{n+1} & F_n \\ F_n & F_{n-1} \end{pmatrix},$$

which is the basis of the algorithm.

From the closed form it follows that the n-th Fibonacci number is given as the upper left element of the matrix

$$\begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}^{n-1}$$

so the problem of computing a Fibonacci number can be reduced to matrix multiplication. For an implementation, it is now important to perform the matrix multiplications efficiently. However, we only have to compute matrix powers, which a computer can do efficiently by using the binary representation of a natural number.

We will explain the procedure for the simpler case of raising a real number to an integral power. Say we want to compute a^d for $a \in \mathbb{R}, d \in \mathbb{N}$. The binary representation of d is given by a finite sequence $\{b_i\}_{i=0}^n, b_i \in \{0, 1\}$ for an implementation dependent $n \in \mathbb{N}$. Now,

$$a^d = \prod_{i=0}^n a^{b_i \cdot 2^i}$$

So, instead of multiplying *a d*-times by itself, we use the following procedure: starting with r = 1, for *i* from 0 to *n*, we multiply *r* with $a^{b_i \cdot 2^i}$. In each iteration, the $a^{(2^i)}$ is computed as $(a^{(2^{i-1})})^2$

The procedure is summerized in the following algorithm:

Algorithm 1 Computation of the n-th Fibonacci number

1: input: natural number n2: set d = n - 1, $A = \begin{pmatrix} 1 & 1 \\ 1 & 0 \end{pmatrix}$, $M = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix}$ 3: while d > 0 do 4: if rightmost-bit of d is 1 then 5: $M = M \cdot A$ 6: end if 7: $A = A^2$ 8: shift d to the right by 1 bit 9: end while 10: output: $F_n = M_{00}$

• Implement this algorithm.

Hints:

- use typedef unsigned long long int Integer to be able to compute larger Fibonacci numbers than in exercise 1 (you should be able to compute F_{200} in the computer pool room).
- write a function void MatMatMult (Integer M[2][2], Integer A[2][2]) for computing the product of two 2 × 2 matrices. The product should be written to M.
- write a function void square (Integer A[2][2]) to compute the square of a matrix. The result should be returned in A.
- checking whether the rightmost-bit of d is 1 can be done by (d & 1) (& is the bitwise and operator).
- shifting d to the right by 1 bit can be done by d = d >> 1; (>> is the shift right operator).
- for unsigned long long int, use %11u in printf statements.