

Computable Analysis

SS 2013, Exercise Sheet #10

EXERCISE 18:

- a) Let P denote a second-order polynomial in $n \in \mathbb{N}$ and in $\lambda \in \mathbb{N}^{\mathbb{N}}$.
Prove by structural induction on P that, for $\lambda \in \mathbb{N}[N]$,
- i) $P(n, \lambda)$ is a polynomial in n
 - ii) whose degree depends only on $\deg(\lambda)$ and
 - iii) that $\mathbb{N} \ni d \mapsto \deg P(n, N^d)$ is a polynomial in d .
- b) Theorem 4.7h) in the script asserts that the closed image of a closed ψ_{\leq}^d -computable set under a computable function is again ψ_{\leq}^d -computable. We also know from Exercise 13c) that intersection of non-empty closed sets is $(\psi_{>}^d \times \psi_{>}^d, \psi_{\leq}^d)$ -computable.
- i) Prove that the image of compact subsets under a computable function is ψ_{\leq}^d -computable.
More precisely use Theorem 4.7j+k) to establish that $C(\mathbb{R}^d, \mathbb{R}^k) \times \mathcal{A}^{(d)} \ni (f, A) \mapsto f[A \cap [0; 1]^d] \in \mathcal{A}^{(k)}$ is $([\rho^d \rightarrow \rho^k] \times \psi_{>}^d, \psi_{\leq}^d)$ -computable.
 - ii) How about $(\psi_{>}, \psi_{>})$ -computability of the mapping $A \mapsto \sin[A]$ for arbitrary compact $A \subseteq \mathbb{R}$?
- c) The lecture constructed a $\psi_{>}$ -computable closed non-empty $A \subseteq [0; 1]$ containing no computable point.
- i) Prove that every $\psi_{>}$ -computable closed non-empty *convex* $C \subseteq [0; 1]$ contains a computable point.
 - ii) Generalize i) to ψ_{\leq}^d -computable closed non-empty convex $C \subseteq [0; 1]^d$.
Hint: induction on d using b i) and Exercise 13c).