

Computable Analysis
SS 2013, Exercise Sheet #6

EXERCISE 14:

- a) Prove that the composition of polynomial-time computable total $f, g : \{0, 1\}^* \rightarrow \{0, 1\}^*$ is again polynomial-time computable.
- b) Compare the computational complexity of the Hamiltonian Cycle problem HC with that of the Hamiltonian Path problem:

$$\text{HP} = \{ \langle G \rangle : G \text{ graph admits a path visiting every vertex precisely once.} \}$$

- c) Construct a polynomial-time reduction from EC (Eulerian Cycle) to HC.
- d) A graph with *integer edge weights* is a mapping $E : V \times V \rightarrow \mathbb{N}$. The *weighted length* of a path (v_1, \dots, v_m) is $\sum_{j=1}^{m-1} E(v_j, v_{j+1})$.
Construct a polynomial-time reduction from HC (Hamilton Cycle) to the following *Traveling Salesperson Problem*:

$$\text{TSP} = \{ \langle \text{bin}(E), \text{bin}(k) \rangle : E \text{ weighted graph admits a Hamiltonian cycle of length } \leq k \}$$

- e) A graph (V, E) is *k-colorable* if it admits a mapping $c : V \rightarrow \{1, \dots, k\}$ with $c(u) \neq c(v)$ for all $(u, v) \in E$. Formalize the problem of whether a given graph is 3-colorable and construct a polynomial-time reduction to SAT.