

>

Background know-how

Limits of Floating-Point arithmetic in C

```
#include <stdio.h>
int main(void) {
    double x=0.7;
    int i = 0;
    while(i < 10) {
        x = 11.0 * x - 7.0;
        printf("%d: %.20lf\n", i, x);
        i=i+1;
    }
}
```

The result of the C-program is rubbish. In the last round it is
 $y = -1127140547773912.5$

Limits of Floating-Point arithmetic in Maple

```
> restart; x :=  $\frac{7.0}{10}$ ;
x := 0.7000000000000000
```

```

> for i from 1 to 30 do
    x := 11·x - 7;
end do;
> x;
0.700000000
(2)

```

> *restart*; $x := \frac{1.0}{3} :$

```

> for i from 1 to 30 do
    x := 3·x - 2;
end do;

```

$$> x; \quad -10294.22328 \quad (3)$$

```
> x := 0 : t := time( ) :  
  for i from 1 to 5000000 do  
    r := rand( ) mod 10;  
    for j from 1 to r do  
      x := x + 1;  
    end do;  
  end do;
```

$x, \text{time}() - t;$

22492822, 91.588

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Numbers, their representations and more and less native number representations for a digital computer

numbers can be elements from various sets. e.g. $x \in \mathbb{Z}$, $x \in \mathbb{N}$.
each number has various representations. e.g.

17

XVII

IIII IIII IIII II

usually, we encode numbers with the help of base-10 digits, i.e. the alphabet
 $\Sigma = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9\}$.

A string $s = (a_{n-1}a_{n-2} \dots a_1a_0) \in \Sigma^n$ is then interpreted as

$$\sum_{i=0}^{n-1} a_i \cdot 10^i : \text{Example: } 17 = 1 \cdot 10^1 + 7 \cdot 10^0$$

What happens, if we use another base, another alphabet?

E.g. with "bits", we have:

$$\Sigma_2 = \{0, 1\} \quad 17_{10} = 1 \cdot 2^4 + 0 \cdot 2^3 + 0 \cdot 2^2 + 0 \cdot 2^1 + 1 \cdot 2^0 = 10001_2$$

$$\Sigma_{16} = \{0, 1, 2, 3, 4, 5, 6, 7, 8, 9, a, b, c, d, e, f\}$$

$$17_{10} = 1 \cdot 16^1 + 1 \cdot 16^0 = 0x11 \quad (\text{so called hex numbers})$$

integer variables of fixed length are the most natural and mostly used kind of variables

Bitstrings are interpreted as numbers in the dual number system.

0 1 0 0 0 1 1 0 0 0 1 0 1 0 1 1 0 1 0 1 0 0 0 0 1 0 0 0 0 1 0 1

.....

.....

bit 0

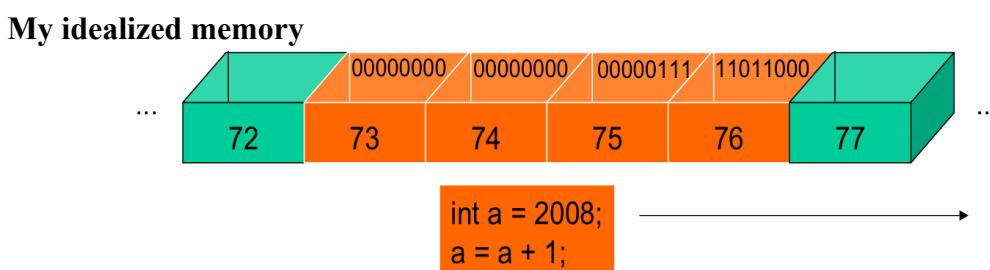
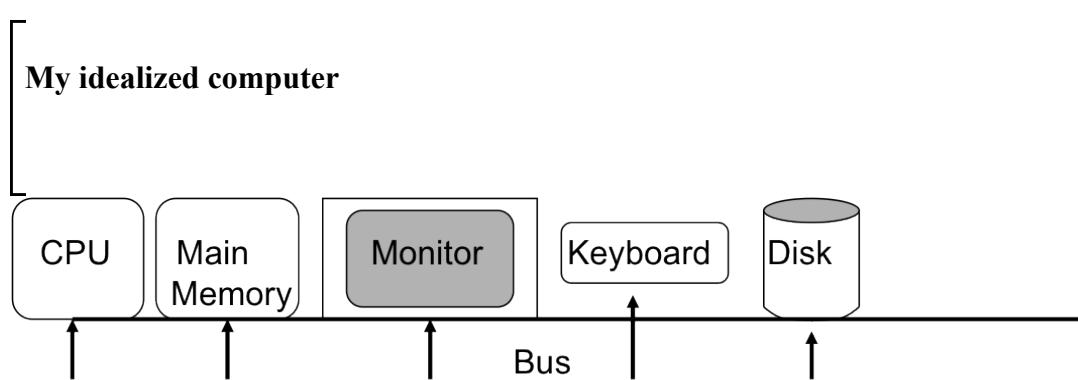
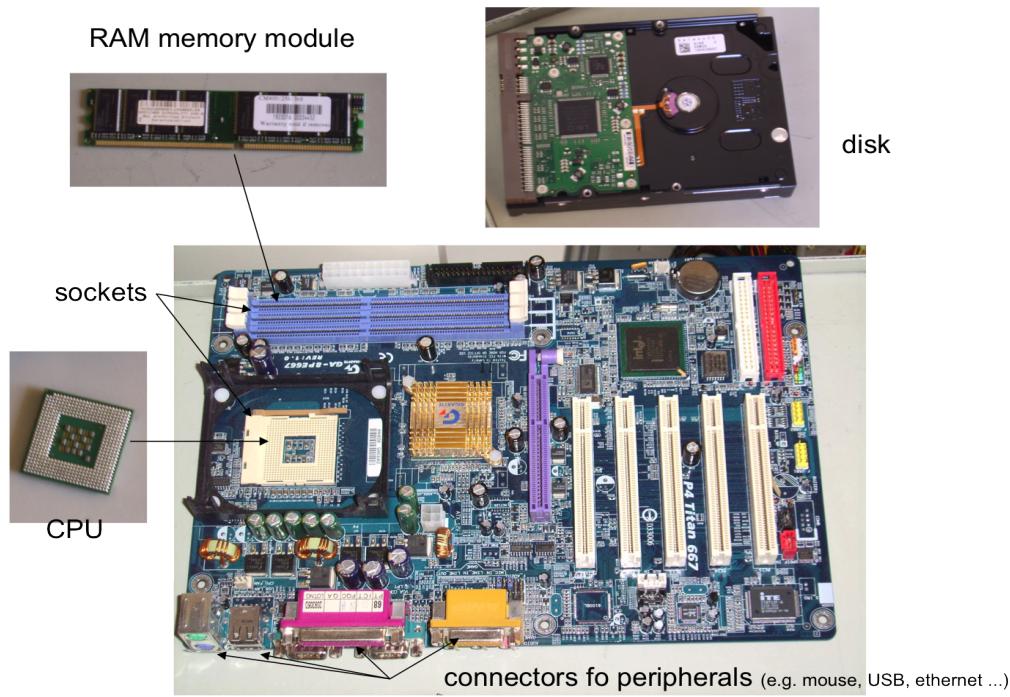
bit 1

.....

bit 30

bit 31 (MSB)

The value then is $bit_{31} \cdot 2^{31} + bit_{30} \cdot 2^{30} + \dots + bit_0 \cdot 2^0$.



How to compute with binary numbers?

□

base-2

$$\begin{array}{r} \text{sum:} \\ + \end{array} \quad \begin{array}{r} 1 & 0 & 1 & 1 \\ \hline 1 & 1 & 1 & 0 \end{array}$$

base-10

$$\begin{array}{r} & 9\ 9 \\ + & \underline{-1-1} \ 3 \\ & 1\ 0\ 2 \square \end{array}$$

1

product:

$$\begin{array}{r} \underline{1011 \cdot 101} \\ 1011 \\ 0000 \\ \hline 1011 \\ 110111 \end{array}$$

Generalized binary fixed-point and floating-point numbers

0.75

$$0.75 = 1 \cdot \frac{1}{2} + 1 \cdot \frac{1}{4} = 0.11_2$$

0.7

$$0.7 = 1 \cdot \frac{1}{2} + 0 \cdot \frac{1}{4} + 1 \cdot \frac{1}{8} + 1 \cdot \frac{1}{16} + \dots$$

the first 64 bits:

0.7 is a periodic number in the binary system.

floating point variables

0/1 sequences are interpreted as sign (s) , mantissa (m) and exponent (p)

0 0 0 0 0 0 0 1 0 1 0 1 0 1 0 1 1 0 1 0 1 0 0 0 0 1 0 0 0 0 1 0 1
s p m

the resulting number then is $s \cdot m \cdot 2^p$

In the IEEE-754 standard, 127 is added to the exponent, and the leading 1 of the mantissa is not stored. The exponent has 8 bits and the mantissa 23 explicit bits, thus 24 implicit bits.

Explicit bits, thus \mathbf{z}' implicit bits.

s p' m'

-> representation errors in IEEE format is not avoidable

-> $x = 0.7$; $x = 11.0 \cdot x - 7.0$; increases the error by a factor of 10

Wrong results in spite of exact computations

Expand

$$\text{expand}\left(\frac{x \cdot (x^3 + 3)}{x \cdot (x + 1)}\right);$$

11379726889978832996251
22492823

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The case

The fibonacci series is defined as follows:

$\text{fib}(0) = 0$, $\text{fib}(1) = 1$ and $\text{fib}(n+1) = \text{fib}(n-1) + \text{fib}(n)$

We would like to know whether $f(n)$ might be expressible as

$$\text{fib}(n) = \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^n - \left(\frac{1 - \sqrt{5}}{2} \right)^n \right)$$

We would like to get some information fast and without lots of hand work.
How can we start working at the exercise? How can Maple help us?

Solution:

Relativly soon, it is clear that:

$$\frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^0 - \left(\frac{1 - \sqrt{5}}{2} \right)^0 \right) \quad 0 \quad (6)$$

and

$$\frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^1 - \left(\frac{1 - \sqrt{5}}{2} \right)^1 \right) \quad 1 \quad (7)$$

Additionally, it must be true that

$$\begin{aligned} & \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^{n-1} - \left(\frac{1 - \sqrt{5}}{2} \right)^{n-1} \right) + \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^n - \left(\frac{1 - \sqrt{5}}{2} \right)^n \right) = \frac{1}{\sqrt{5}} \\ & \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^{n+1} - \left(\frac{1 - \sqrt{5}}{2} \right)^{n+1} \right) \end{aligned}$$

Some large numbers can quickly be tested, the expression may be simplified via the command `simplify`. An example is 876:

$$simplify \left(\frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^{876-1} - \left(\frac{1-\sqrt{5}}{2} \right)^{876-1} \right) + \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^{876} - \left(\frac{1-\sqrt{5}}{2} \right)^{876} \right) - \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^{876+1} - \left(\frac{1-\sqrt{5}}{2} \right)^{876+1} \right) \right) \quad (8)$$

The procedure becomes by far more tricky, if we want Maple to show equality for general n . Sometimes, it helps to expand the expression.

$$\begin{aligned}
& \text{expand} \left(\frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^{n-1} - \left(\frac{1-\sqrt{5}}{2} \right)^{n-1} \right) + \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^n - \left(\frac{1-\sqrt{5}}{2} \right)^n \right) \right. \\
& \quad \left. - \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^{n+1} - \left(\frac{1-\sqrt{5}}{2} \right)^{n+1} \right) \right); \\
& \frac{1}{5} \frac{\sqrt{5} \left(\frac{1}{2} + \frac{1}{2} \sqrt{5} \right)^n}{\frac{1}{2} + \frac{1}{2} \sqrt{5}} - \frac{1}{5} \frac{\sqrt{5} \left(\frac{1}{2} - \frac{1}{2} \sqrt{5} \right)^n}{\frac{1}{2} - \frac{1}{2} \sqrt{5}} + \frac{1}{10} \sqrt{5} \left(\frac{1}{2} + \frac{1}{2} \sqrt{5} \right)^n \\
& \quad - \frac{1}{10} \sqrt{5} \left(\frac{1}{2} - \frac{1}{2} \sqrt{5} \right)^n - \frac{1}{2} \left(\frac{1}{2} + \frac{1}{2} \sqrt{5} \right)^n - \frac{1}{2} \left(\frac{1}{2} - \frac{1}{2} \sqrt{5} \right)^n \tag{9}
\end{aligned}$$

$$h := \text{simplify} \left(\frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^{n-1} - \left(\frac{1-\sqrt{5}}{2} \right)^{n-1} \right) + \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^n - \left(\frac{1-\sqrt{5}}{2} \right)^n \right) - \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1+\sqrt{5}}{2} \right)^{n+1} - \left(\frac{1-\sqrt{5}}{2} \right)^{n+1} \right); \right. \\ \left. - \frac{1}{10} \sqrt{5} \left(- \left(\frac{1}{2} + \frac{1}{2} \sqrt{5} \right)^n - \sqrt{5} \left(\frac{1}{2} + \frac{1}{2} \sqrt{5} \right)^n + \left(\frac{1}{2} - \frac{1}{2} \sqrt{5} \right)^n - \sqrt{5} \left(\frac{1}{2} - \frac{1}{2} \sqrt{5} \right)^n + 2 \left(\frac{1}{2} + \frac{1}{2} \sqrt{5} \right)^{n+1} - 2 \left(\frac{1}{2} - \frac{1}{2} \sqrt{5} \right)^{n+1} \right) \right); \quad (11)$$

$$is \left(simplify \left(\left(\frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^n - \left(\frac{1 - \sqrt{5}}{2} \right)^n \right) + \frac{1}{\sqrt{5}} \cdot \left(\left(\frac{1 + \sqrt{5}}{2} \right)^n \right. \right. \right. \right. \right. \\$$

$$-\left(\frac{1-\sqrt{5}}{2}\right)^n\right)\right)-\left(\frac{1}{\sqrt{5}}\cdot\left(\left(\frac{1+\sqrt{5}}{2}\right)^{n+1}-\left(\frac{1-\sqrt{5}}{2}\right)^{n+1}\right)\right)=0;$$

false

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```

>
>
>
> restart; sff := [seq(0, i=0..99)]:
> sff[0] := 0; sff[1] := 1;
Error, out of bound assignment to a list
sff1 := 1

```

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> sff[0 + 1] := 0; sff[1 + 1] := 1;
sff1 := 0
sff2 := 1

```

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> sff[2 + 1] := sff[0 + 1] + sff[1 + 1];
sff3 := 1

```

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```

> for i from 3 to 99 do
  sff[i + 1] := sff[i - 1 + 1] + sff[i - 2 + 1];
  if i = 97 then print(sff[i + 1]) fi;
end do:
83621143489848422977

```

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```

> restart; sff := Array(0 .. 10000, fill = 0):
>
> sff[0] := 0; sff[1] := 1;
sff0 := 0
sff1 := 1

```

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```

> sff[2] := sff[0] + sff[1];
sff2 := 1

```

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```

> for i from 3 to 10000 do
  sff[i] := sff[i - 1] + sff[i - 2];
  if i = 97 then print(sff[i]) fi;
end do:
83621143489848422977

```

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```

> sff1 := sff[0]; sff2 := sff[1]; sff3 := sff[2];
sff1 := 0
sff2 := 1
sff3 := 1

```

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```

> for i from 3 to 10000 do
  sff4 := sff3 + sff2;
  sff2 := sff3 : sff3 := sff4;
  if i = 97 then print(sff3) fi;
end do:

```

[>

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