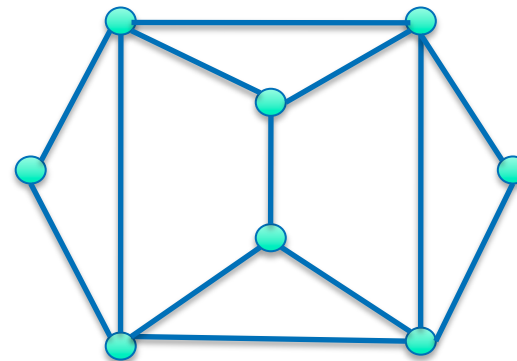
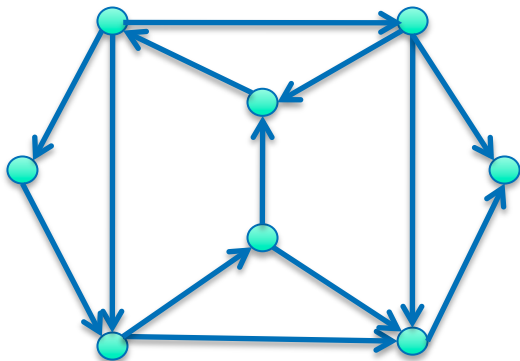


Graphs and classic graph problems

- What is a graph?
 - An undirected **graph** is a pair $G=(V,E)$, with $E \subseteq \{\{u,v\} \mid u,v \in V\}$.
 - The elements of E are not ordered.
 - Elements of V are called **nodes**, or **vertices resp.**
Elements of E are called **edges**
 - A **directed graph** is a pair $G=(V,E)$ as well. However, the elements of E are ordered pairs of elements of V . Thus
 $E \subseteq \{(u,v) \mid u,v \in V\}$.
 - Elements of V are called nodes
Elements of E are called **directed edges** (in ger.: **gerichtete Kanten** or **Bögen**)



Graphs and classic graph problems

Neighborhood relations

Incidence: A node v is said to be incident to an edge e if $v \in e$.



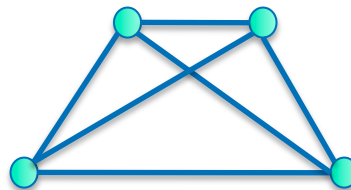
adjacency: Two nodes are called adjacent in G , if $\{x,y\} \in E$.



Degree: The degree of a node v [$\deg(v)$] is the number of with v incident edges.

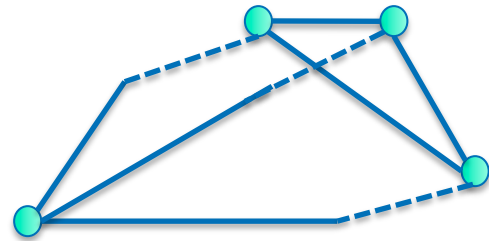
Clique: If any two nodes of G are neighbors, G is called the complete graph.

G is also called a clique.



Graphs and classic graph problems

Claim: $\sum_{v \in V} \deg(v) = 2|E|$



Proof: counting edges

Claim: $\{v \mid \deg(v) \text{ odd}\}$ is even

Proof: clear

Graphs and classic graph problems

■ Paths and Cycles

- A **path** in a graph (directed or non-directed) $G = (V, E)$ is a sequence of nodes (v_0, v_1, \dots, v_k) from V , such that for all $i \in \{1, \dots, k\}$ is valid that $(v_{i-1}, v_i) \in E$ (or $\{v_{i-1}, v_i\} \in E$, resp.) is an edge.
If $i < j$, v_i is called a predecessor of v_j , and v_j a successor of v_i
- A path is called **simple**, if no node occurs more than once in the sequence (v_1, \dots, v_k) .
- A simple path is called **cycle**, if $v_0 = v_k$.
- A graph is called **acyclic** (ger.: kreisfrei), if there is no cycle contained in the graph.

Graphs and classic graph problems

■ Connected components

- A undirected graph G is **connected**, if there is a path from v to w for all node pairs $v, w \in V$.
- A directed graph is strongly connected, if there are paths from all nodes to any other node.
- Connected parts of G are called **connected components**.

Reminder: A set V with binary relation is called an equivalence relation if for all for all $x, y, z \in V$:

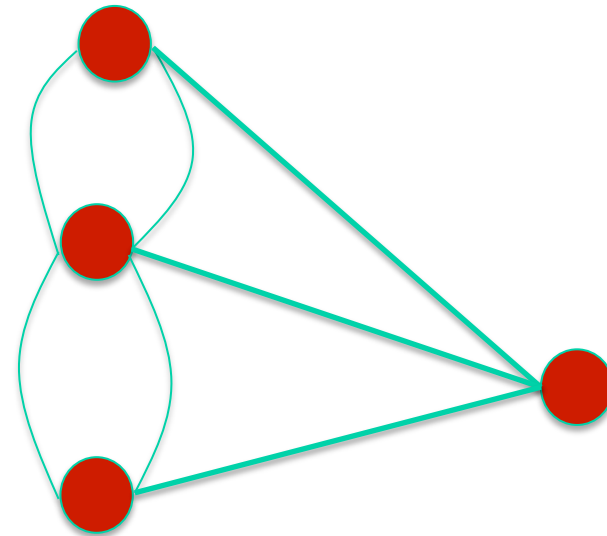
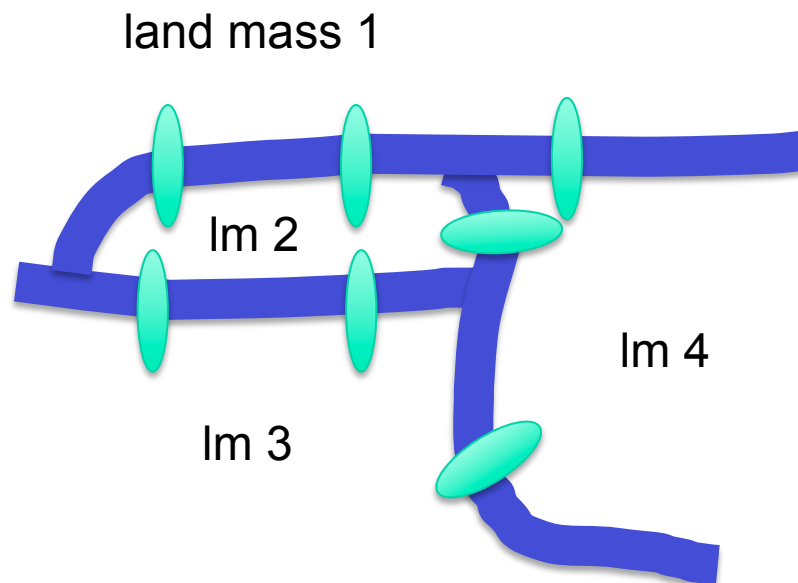
- $x \sim x$, reflexive
- $x \sim y \Rightarrow y \sim x$, symmetric
- $x \sim y$ and $y \sim z \Rightarrow x \sim z$, transitive

Observations: Connected components of a graph form equivalence classes.

Graphs and classic graph problems

Seven bridges of Königsberg (Euler 1736):

Is it possible was to find a walk through the city that would cross each bridge once and only once? Such a walk is called *Eulerian path*.

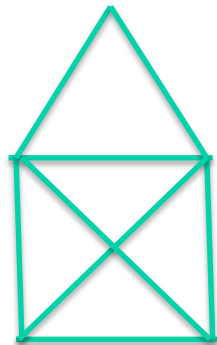


An Eulerian path with identical start and end point is called *Eulerian cycle*.
→ no, both do not exist in Königsberg of 1736.

Graphs and classic graph problems

Seven bridges of Königsberg (Euler 1736):

Claim: A necessary and sufficient condition for a walk of the desired form is that the graph is connected and has exactly zero (-> Eulerian cycle) or two nodes (-> Eulerian path) of odd degree.



Haus vom Nikolaus contains

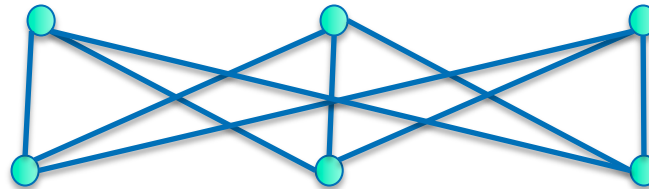
- 1 node of degree 2
- 2 nodes of degree 4
- 2 nodes of degree 3

→ Contains an Eulerian path, but no Eulerian cycle.

Graphs and classic graph problems

Bipartite graphs

A graph is called bipartite, if it is possible to partition V into two disjoint parts V_1 and V_2 , such that each edge has one end in V_1 and the other one in V_2 .

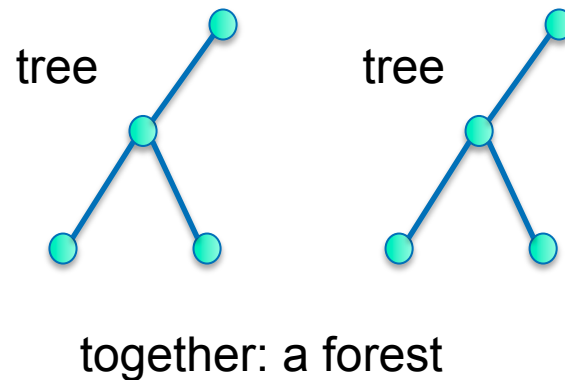


Claim: A graph is bipartite if and only if any cycle in the graph has even length.

Proof: Exercise

Graphs and classic graph problems

- A **tree** is a connected graph without any cycles
- A **forest** is a graph without any cycles, thus a collection of trees.
- Nodes with degree 1 are called **leaf**, the others **inner nodes**.



Claim: A forest F has $|E| = |V| - C(F)$ edges, with C number of components.

Claim: If $|E| > |V| - C(F)$, G will contain a cycle.

Claim: For any tree is valid: $|E| = |V| - 1$

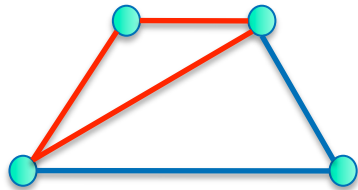
Proof: Exercise

Graphs and classic graph problems

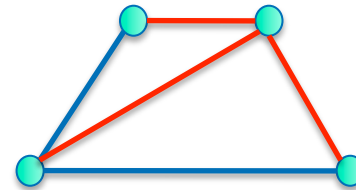
- subgraphs, spanning subgraphs

- $G'=(V',E')$ is called **subgraph** of $G=(V,E)$, if $V' \subseteq V$ and $E' \subseteq E$.

- $G' \subseteq G$ is called a **spanning subgraph** of G , when additionally: $V'=V$



Graph with subgraph



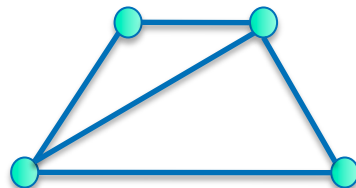
spanning tree

Claim: A graph is connected iff it contains a spanning tree.
(Proof: \Leftarrow : clear; \Rightarrow : remove edges from cycles)

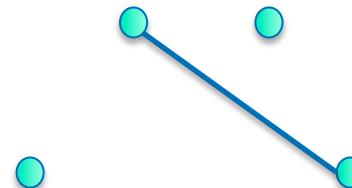
Graphs and classic graph problems

■ Complementary Graph

- The for G **complementary graph** G' is the graph $G'=(V,E')$, with $(i,j) \in E' \Leftrightarrow (i,j) \notin E$



graph



complement

- Claim: At least one of the graphs G or G' is connected.
Proof: Exercise.